# Networks out of Control:

# Homework Set 2

## Exercise 1

We identify a weaker condition for G(n, p) to have diameter at most 2, than the one seen in class: Show if  $p \ge c \frac{\log n}{\sqrt{n}}$ , where c > 0 is some constant, then G(n, p) has diameter at most 2 a.a.s.

# Exercise 2

Consider two families of random graphs  $G_n = G(n, p)$ , with  $p(n) = n^{-6/5}$ , and another independent random graph  $H_n = G(\log n, (\log n)^{-8/7})$ . Can you show that  $H_n$  does not appear in (i.e., is not a subgraph of)  $G_n$  for large n?

Hint: Identify a subgraph that is likely to appear in H, then show that this subgraph does not appear in G.

#### Exercise 3

We have seen in class that t(n) = 1/n is a threshold function for the appearance of cycles of *fixed* order k in G(n, p). Observe that this does not directly imply that t(n) is also a threshold function for the appearance of cycles of *any* order. Prove that the assertion is in fact true.

## Exercise 4

Let V denote, as usual, the vertex set of G(n, p), with 0 constant. A subset W of V is called*independent*if and only if no pair of vertices of W are adjacent. Show that <math>G(n, p) contains a.a.s no independent set of more than  $2c \log_q n$  vertices, where c > 1 and q = 1/(1-p).

Hint: Let  $X_k$  denote the number of independent subsets of V, with k vertices. Compute first  $\mathbb{E}[X_k]$ , and next an upper bound of  $\mathbb{E}[X_k]$  with  $k = 2c \log_q n$ .