

A note about the analysis of MH for coloring.

Recall we have a chain in space of proper colorings:

- 1) Select $v \in V$ at random;
- 2) Select $c \in \{1 \dots q\}$ and recolor v if c is allowed.

LEMMA:

We want to show this chain is irreducible if $q \geq \Delta + 2$.

$$\Delta = \max_{v \in V} \deg(v)$$

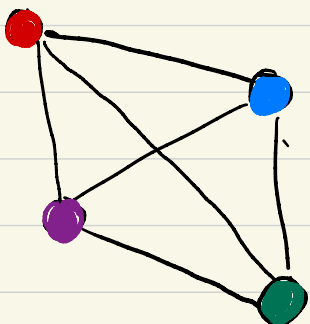
It should be clear that it is enough to show that any two assignments \underline{x} and \underline{y} can be connected by a path

$$\underline{x} \rightarrow \underline{z}_1 \rightarrow \underline{z}_2 \rightarrow \dots \rightarrow \underline{z}_m \rightarrow \underline{y}$$

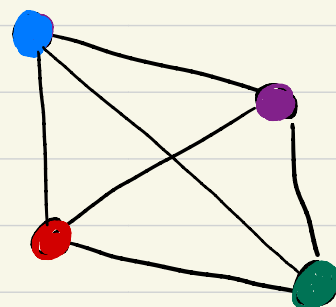
such that two successive assignments differ by only one color:

Question: Consider the following working example and find an algorithm assuming $q \geq \Delta + 2$.

assignment \underline{x}



assignment \underline{y}



$$\Delta = 4$$

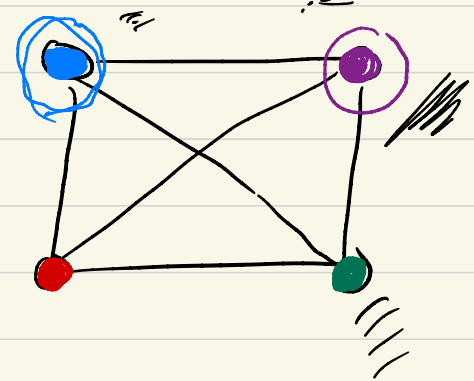
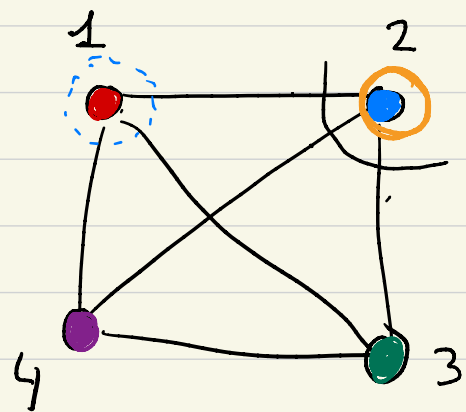
$$c \in \{ \text{red, blue, green, purple, orange} \}$$

$\Delta + 2$ colors.

Available colors {      }


target:    

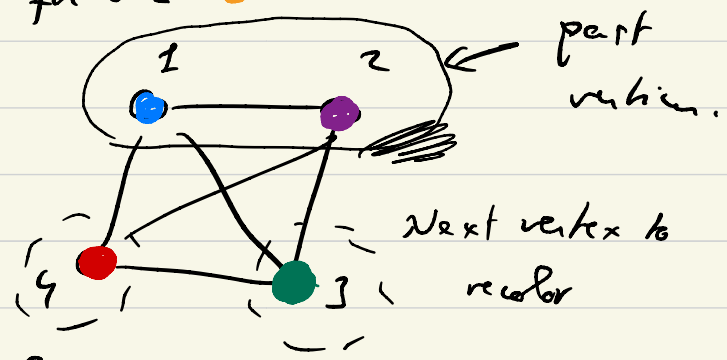
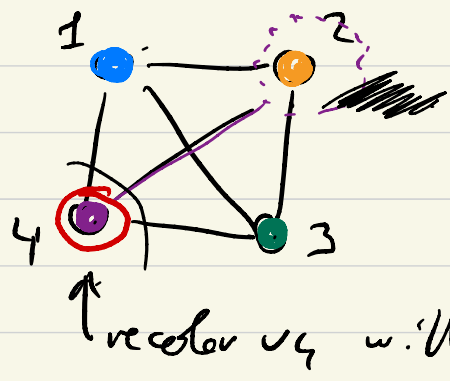
initial assignment X



Sequence of recolorings: $X \rightarrow z_1 \rightarrow z_2 \rightarrow \dots \rightarrow Y$

v_1 is recolored in blue Trial \rightarrow does not work.

First I take a provisional color for v_2 



DONE!

Main Idea of Algorithm:

- Order vertices and sequentially take vertex in that order and try to change their color to match those of Y .
- If you find a conflict (with future vertex) then before coloring current vertex you recolor the conflicting in no future with an available color. $q \geq \Delta + 2$ ensures \exists this available color.