| First | name: |
|---------|-------|
| 1 11 31 | name. |

Family name:

EXAM TCP/IP NETWORKING Duration: 3 hours

Jean-Yves Le Boudec

January 2020

INSTRUCTIONS

- 1. Write your solution into this document and return it to us (you do not need to return the figure sheet). You may use additional sheets if needed. Do not forget to write your name on **each of the four problem sheets** and **all** additional sheets of your solution.
- 2. All problems have the same weight.
- 3. You may need to make additional assumptions in order to solve some of the questions. If this happens, please describe such assumptions explicitly.
- 4. Figures are on a separate sheet, for your convenience.
- 5. No documents, no electronic equipments are allowed, except for a non-connected calculator.

PROBLEM 1

Consider the network for problem 1 in the figure sheet. A, B, C, D and E are hosts; BR1, BR2 and BR3 are bridges running the spanning tree protocol; N is an IPv4 NAT; R1, R2 and R3 are routers. Plain lines are physical connections. O1 to O6 are points where we observe traffic.

Host A is IPv6-only; C and E are IPv4-only; B and D are dual stack. All routers are dual-stack.

MAC addresses are denoted with e.g. D, E, BR3e, ..., Nn. If you need to make assumptions about addresses, please write them explicitly.

All links are full duplex Ethernet. We assume that all machines are correctly configured (unless otherwise specified), proxy ARP is not used and there is no VLAN.

- 1. ISP4 allocates to *N*-north the address 11.12.13.14.
 - (a) Give possible values of x, y, z (in the IPv4 addresses of B, C, D) and the network masks at N-west, B and C.

(b) D sends one UDP datagram to E. We observe the resulting packet at observation point O5. What are the MAC source and destination addresses ?

(c) B sends a sequence of IP packets to D. We assume all routing and bridging protocols have converged and all forwarding tables have been learnt. At which of the observation points O2, O3, O4 are the packets visible ? Justify your answer.

(d) C downloads a huge file from a web server at E using HTTP over QUIC. C uses the local port 4567. The server port number for QUIC is 443. At the same time, B also downloads a file from E using HTTP over QUIC. By coincidence, B uses the same local port number, namely 4567. We observe the packet headers in the packets resulting from this transfer at O5 and O6, in the direction to E. Give possible values of the protocol, the source and destination port numbers and the source and destination IP addresses. Give the answers in the tables below.

| At observation point O5, towards E: | | | | | |
|-------------------------------------|--------------------------|-----------------------------------|---|--|--|
| IP dest | protocol | source port | dest. port | | |
| | | | | | |
| | | | | | |
| | | | | | |
| | | | | | |
| | 5, towards E: IP dest | 5, towards E: IP dest protocol | 5, towards E: IP dest protocol source port | | |

| At observation point O6, towards E: | | | | |
|-------------------------------------|---------|----------|-------------|------------|
| IP source | IP dest | protocol | source port | dest. port |
| | | | | |
| | | | | |
| | | | | |
| | | | | |

- 2. ISP6 delegates the prefix 2001:a:a:a::/76 to the network of Problem 1. Inside the network, all IPv6 subnet prefixes are /80.
 - (a) Give the uncompressed version of the address 2001:a:a:a:1::2.

(b) Among the following addresses, which ones are possible for host B? Put an X in the correct

boxes in the table below, with a short justification.

| address | possible | not possible |
|------------------|----------|--------------|
| 2001::3 | | |
| 2001:a:a:a:2::3 | | |
| 2001:a:a:a:5::3 | | |
| 2001:a:a:f::3 | | |
| 2001:a:a:a:11::3 | | |
| Justification: | | |
| | | |
| | | |
| | | |
| | | |

(c) B downloads a huge file from a web server at A using HTTP over TLS over TCP. B uses the local port 4567 and the server port at A is 443. At the same time, D also downloads a file from A, also using HTTP over TLS over TCP. By coincidence, D uses the same local port number, namely 4567. We observe the IP headers in the packets resulting from this transfer at O1, in the direction towards A. Give possible values of the protocol, the source and destination port numbers and the source and destination IP addresses. Give the answers in the tables below.

| At observation point O1, towards A: | | | | |
|-------------------------------------|---------|------|-----------|------------|
| IP source | IP dest | prot | src. port | dest. port |
| | | | | |
| | | | | |
| | | | | |
| | | | | |

| First name: | Family name: | |
|-------------|--------------|--|

PROBLEM 2

Consider the network for problem 2 in the figure sheet. There are three ASs, A, B and C with routers A1, A2, B1, B2, B3, B4, R1, R2, C1 and C2. The physical links are shown with plain lines. Each AS uses OSPF with Equal Cost Multipath as IGP, and every router inside each AS uses OSPF. The cost of every link and every directly attached network is 1, except when otherwise shown on the figure.

The figure shows stub networks, at routers A1, C2, R1 and R2, with their IPv6 address prefixes. The lower case symbols such as b1w, a1e also represent IPv6 addresses.

Routers A1, A2, B1, B2, B3, B4, C1 and C2 use BGP with their external neighbours and as required with their internal neighbours. The routers R1 and R2 may or may not use BGP, depending on the question. No confederation or route reflector is used.

We assume that the BGP decision process use the following criteria in decreasing order of priority. BGP identifiers are router names such as A1, A2...

- 1. Shortest AS-PATH
- 2. E-BGP is preferred over I-BGP
- 3. Shortest path to NEXT-HOP, according to IGP
- 4. Lowest BGP identifier is preferred (e.g. A1 is preferred over A2)

Furthermore, we assume that:

- No optional BGP attribute (such as MED, LOCAL-PREF etc.) is used in any BGP message.
- No aggregation of route prefixes is performed by BGP.
- The policy in A, B, C is such that all available routes are accepted and propagated to neighbouring ASs, as long as the rules of BGP allow.
- Every router redistributes internal OSPF destinations into BGP.
- Every router performs recursive forwarding-table lookup.
- Equal Cost Multi-Path routing is supported by all routers.

- 1. In this question, we assume that R1 and R2 run BGP. At time t_1 , BGP and OSPF have converged in all ASs.
 - (a) At time t_1 , what is the list of BGP routes received by A2 with destination = 2001:1::/32 ? Which route is selected as best route by A2 ? Give your answer in the table below, with a short justification (put as many rows as necessary).

| At <i>A</i> 2 : | | | | |
|-----------------|---------------------|--------------|---------|--------------|
| From BGP Peer | Destination Network | BGP Next-Hop | AS-Path | Best route ? |
| | 2001:1::/32 | | | |
| | 2001:1::/32 | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| Justification: | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |

(b) Still at time t_1 , what is the list of BGP routes received by A2 with destination = 4001:1::/32 ? Which route is selected as best route by A2 ? Give your answer in the table below, with a short justification (put as many rows as necessary).

| At A2 : | | | | |
|----------------|---------------------|--------------|---------|--------------|
| From BGP Peer | Destination Network | BGP Next-Hop | AS-Path | Best route ? |
| | 4001:1::/32 | | | |
| | 4001:1::/32 | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| Justification: | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |

(c) Still at time t_1 , what is the list of BGP routes received by R_2 with destination = 4001:1::/32 ? Which route is selected as best route by R_2 ? Give your answer in the table below, with a short justification (put as many rows as necessary).

| At <i>R</i> 2 : | | | | |
|-----------------|---------------------|--------------|---------|--------------|
| From BGP Peer | Destination Network | BGP Next-Hop | AS-Path | Best route ? |
| | 4001:1::/32 | | | |
| | 4001:1::/32 | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| Justification: | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |

(d) Still at time t_1 , R_2 has a packet to forward with destination address 4001:1:2:3::1. Which path will this packet take inside AS B? At which routers along this path is recursive table lookup required for forwarding this packet?

(e) At time $t_2 > t_1$ the link A1 - C1 breaks. At time $t_3 > t_2$, the BGP protocol has converged again. what is the list of *valid* BGP routes received by A2 with destination = 4001:1::/32 ? Which route is selected as best route by A2? Give your answer in the table below, with a short justification (put as many rows as necessary).

| At A2 : | | | | |
|----------------|---------------------|--------------|---------|--------------|
| From BGP Peer | Destination Network | BGP Next-Hop | AS-Path | Best route ? |
| | 4001:1::/32 | | | |
| | 4001:1::/32 | | | |
| | | | | |
| | | | | |
| | | | | |
| Justification: | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |

(f) At time $t_3 > t_2$, the link A1 - C1 is repaired and all routing protocols have converged again. Router C1 is compromised and sends the bogus route

dest = 3001:1:2::/48, as-path = C, next-hop = cle

to B4. No other bogus message is sent by any other router. The bogus message is accepted by B4. At time $t_4 > t_3$, BGP has converged again and B2 has a packet to forward with destination address 3001:1:2:3::1. By which link will this packet exit AS B?

- 2. In this question we assume that the network is restarted, with the following changes in the configurations of R1:
 - R1 does not run BGP (but continues to run OSPF).
 - In addition, R1 is configured with a static default route to B2. This route co-exists in the forwarding table with the routes obtained from OSPF.

We also assume that there is no more bogus announcement. Recall that there is no redistribution of BGP into OSPF (but there *is* redistribution of internal OSPF destinations into BGP).

At time t_5 the routing protocols have converged.

R1 has a packet to forward with destination address 4001:1:1:3:3::1. Which path will this packet take inside AS B? At which routers along this path is recursive table lookup required for forwarding this packet?

- 3. In this question we assume that the network is restarted, with the following changes in the configurations of routers inside AS *B*:
 - *R*1 and *R*2 do not run BGP (but continue to run OSPF).
 - There is no static default route in *R*1.
 - Routers B1, B2, B3 and B4 announce in OSPF the destination :: /0 with cost = 5. This is treated by OSPF as an internal destination.

We also assume that there is no more bogus announcement. Recall that there is no redistribution of BGP into OSPF (but there *is* redistribution of internal OSPF destinations into BGP).

At time t_6 the routing protocols have converged. R1 has a large number of packets to forward with destination addresses 4001:1:x where x is a string of 96 bits. The value of x is different for every packet. Which way will these packets travel inside AS B?

| First name: | Family name: |
|-------------|--------------|

PROBLEM 3

Consider the network for problem 3 on the figure sheet.

- Hosts A, B and C are downloading content from server S. R1, R2 and R3 are routers, unless otherwise specified.
- The link rates are indicated on the figure. All links are full duplex with same rate in both directions.
- There is no other system than shown on the figure, and we neglect all flows other than between A, B, C and S. There is no other capacity constraint than the link capacities shown on the figure. We also neglect the impact of the acknowledgement flows in the reverse direction.
- The round trip times (RTTs) shown on the figure are between S and respectively A, B, C (for example, the RTT between S and B is 15 ms. The RTTs include all processing times and all queuing delays at routers.
- We neglect all overheads and assume that the link capacities can be fully utilized at bottlenecks.
- The MSS is the same for all flows and is equal to 1250 Bytes = 10^4 bits.
- We call x, y, z the rates of flows $S \to A, S \to B, S \to C$ in Mb/s.
- 1. What allocations (x, y, z) are Pareto-efficient ?

2. Assume that some bandwidth manager is used, which allocates rates to flows according to max-min fairness. What are the values of x, y, z?

3. Same question with proportional fairness instead of max-min fairness.

4. We now assume that the three flows are using TCP RENO with ECN. What is the value of the rate of each flow ?

5. We continue to assume that the flows are using TCP with ECN. We observe the IP headers of packets on the link from R1 to A. Which proportion of packets do we see marked as "Congestion Experienced"?

6. Assume now that R3 is an application layer gateway instead of a router, namely, the flow sent by S to C is relayed at the application level by R3 to C. We assume that all flows are using TCP with ECN. The round trip time from S to R3 is 15 msec; from R3 to C it is 80 msec; the other round trip times are unchanged. We assume that the application layer gateway has infinite processing power and storage. What is the value of the rate at which C receives data from S via R3 ?

| First name: | Family name: | |
|-------------|--------------|--|

PROBLEM 4

- 1. In a smart grid, a sensor S sends measurements every 20 msec to two data concentrators D1 and D2 (see figure). The sensor uses IP multicast, with source specific multicast. It uses the multicast address ff35::4:3:2:1. The data stream is unidirectional, from the sensor. The network is a single bridged LAN.
 - (a) We do a packet capture at the networking interface of the sensor S and at networking interface of D1. We observe only the packets that carry the unidirectional data stream sent by S. Which addresses do we see in the packets sent by the sensor ? Put the answer in the table below (see on the figure sheet for device addresses).

| At sensor S: | | | | | |
|--------------------------|----------|-----------|---------|--|--|
| MAC source | MAC dest | IP source | IP dest | | |
| | | | | | |
| | | | | | |
| At Data Concentrator D1: | | | | | |
| MAC source | MAC dest | IP source | IP dest | | |
| | | | | | |
| | | | | | |

(b) We want that another machine (SCADA) also receives the measurements sent by the sensor. What is required for that at the sensor S and at the SCADA ?

2. Below is the python code of an application, one proposed by Homer and one proposed by Bart

```
# HOMER, Jan 2020, EPFL
import socket
HOST = 'localhost'
PORT = 5002
sock = socket.socket(socket.AF_INET6, socket.SOCK_STREAM)
sock.setsockopt(socket.SOL_SOCKET, socket.SO_REUSEADDR, 1)
sock.bind((HOST, PORT))
sock.listen(1)
while True:
    connection, addr = sock.accept()
    while True:
       data = sock.recv(16).decode()
       print("received:", data)
       if data != b'':
           sock.sendall(data.encode())
       else:
           print("No more data from", addr)
           break
    connection.close()
# BART, Jan 2020, EPFL
import socket
HOST = 'localhost'
PORT = 5002
sock = socket.socket(socket.AF_INET6, socket.SOCK_STREAM)
sock.setsockopt(socket.SOL_SOCKET, socket.SO_REUSEADDR, 1)
sock.bind((HOST, PORT))
sock.listen(1)
while True:
   connection, addr = sock.accept()
    while True:
       data = connection.recv(16).decode()
       print("received:", data)
       if data != b'':
           connection.sendall(data.encode())
       else:
           print("No more data from", addr)
           break
    connection.close()
```

(a) One of these two programs works, the other causes a run-time error. Which is the correct program ? Justify your answer.

- (b) Say what is true about the version of the application that works (there is exactly one correct answer¹).
 - i. \Box it is a UDP server
 - ii. \Box it is a UDP client
 - iii. \Box it is a TCP server
 - iv. \Box it is a TCP client

¹For question 2 (b), 100% of the points of the question are obtained if only the correct answer is selected; if only one incorrect answer is selected, the grade is negative and is -33%; if zero answer or more than one answer is selected the grade is 0.

3. Say what is true about IP fragmentation (put true/false in the cells below).

| | Host may fragment | Router may fragment | Router may re-assemble |
|-----------|-------------------|---------------------|------------------------|
| With IPv4 | | | |
| With IPv6 | | | |

4. Both H1 and H2 are IPv6-only hosts (see figure). They communicate via IPv6, however H1's local router A does not have native IPv6 access on its wide-area side. Instead, A receives IPv4 public access from an IPv4 provider and uses a tunnel broker offered by IPv6 provider P. P delegates to H1's local area network the prefix 2001:1:2:3::/64; A is the tunnel client and B is the tunnel server. The IPv6 address at A's end of the tunnel is 2001:1:a:b::2 and the IPv6 address at B's end of the tunnel is 2001:1:a:b::1. The IPv6 addresses of H1 and H2 are shown on the figure.

H1 sends one UDP message to H2. The message is small and fits in one IP packet even after encapsulation. We observe the IP packet resulting from this activity at observation points 1, 2 and 3. Give the IP addresses and protocol / next header in the following table.

| At observation point 1, towards H2: | | | | |
|-------------------------------------|-------------|-------------|--|--|
| IPv6 source | IPv6 dest | Next Header | | |
| | | | | |
| At observation point 2, | towards H2: | | | |
| IPv4 source | IPv4 dest | Protocol | | |
| | | | | |
| At observation point 3, | towards H2: | | | |
| IPv6 source | IPv6 dest | Next Header | | |
| | | | | |

TCP IP EXAM - FIGURES

For your convenience, you can separate this sheet from the main document. Do not write your solution on this sheet, use only the main document. You do not need to return this sheet.















Problem 4, Question 1.



