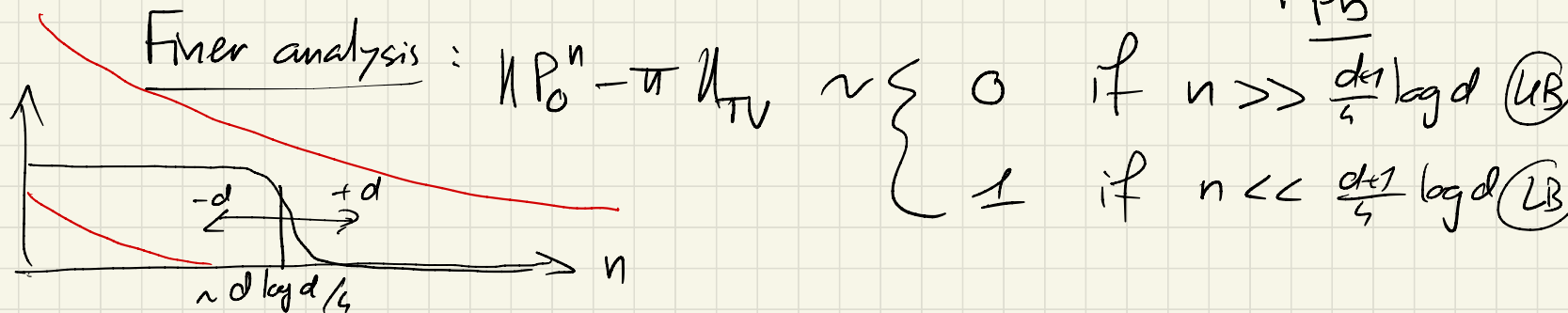


MCAA lecture 8

Recap: cut off phenomenon

1. RW on $S = \{0, 1\}^d$ S very large
(prob $\frac{1}{d+1}$ to stay or to flip one bit)

Bounds: $\frac{1}{2} \exp(-\frac{2n}{d+1}) \leq \|P_0^n - \pi\|_{TV} \leq \frac{1}{2\sqrt{d_0}} \cdot \exp(-\frac{2n}{d+1})$



Pf idea:

all the ev.

$$\underline{UB}: \|P_0^n - \pi\|_{TV} \leq \frac{1}{2} \left(\sum_{z=0}^{2^n} 1_z \right)^{1/2} \dots$$

~ 0 si $n \gg \frac{d \log d}{4}$

$$\underline{LB}: \|P_0^n - \pi\|_{TV} = \max_{A \subset S} |P_0^n(A) - \pi(A)|$$
$$\geq |P_0^n(A) - \pi(A)| \quad \forall A \subset S$$

$$\text{Pick } A_\beta = \left\{ x \in S : \left| |x| - \frac{d}{2} \right| \leq \frac{\beta}{2} \sqrt{d} \right\}$$

$$\Rightarrow \pi(A_\beta) \approx 1, \quad P_0^n(A_\beta) \approx 0 \text{ if } n \ll \frac{d \log d}{4}$$

Card shuffling

$S = \{ \text{permutations of } \{1..N\} \}$ $|S| = N!$
large!

Shuffling method = Markov chain on S

Question: For a given method, how long does
it take to decently shuffle the deck?

(i.e. to have a distribution ϵ -close in TV-dist
to the uniform distribution on S)

Method 0: cut repeatedly the deck

not an ergodic chain ($51!$ equivalence classes)

Method 1: "random to top"

- choose a number unif at random in $\{1..N\}$
- look for the card with this number & put it on top

This is an ergodic chain!

Claim: $\Theta(N \log N)$ shuffles are needed with this method.

Pf: UB

Coupling

X_n = chain starting from the identity state Id

Y_n = chain starting from $\pi \sim$ uniform

choose a number $\in \{1, \dots, N\}$ unif. at random

& look after the card with this number in

each deck, and put it on top in each deck.

Former lemma:

$$\| \underbrace{P_{\text{Id}}^n}_{\text{dist of } X_n} - \underbrace{\pi}_{\text{dist of } Y_n} \|_{\text{TV}} \leq \mathbb{P}(X_n \neq Y_n)$$

Observation:

After each card number has been picked at least once, the two decks are the same

So if $\tau = \inf\{n \geq 1: \text{each number has been picked once}\}$

Then $\mathbb{P}(X_n \neq Y_n) \leq \mathbb{P}(\tau > n)$

Coupon collector pb:

$$E(T) = \frac{N}{N} + \frac{N}{N-1} + \frac{N}{N-2} + \dots + \frac{N}{3} + \frac{N}{2} + \frac{N}{1}$$

$$= N \cdot \sum_{k=1}^N \frac{1}{k} \approx N \cdot \log N$$

$$\text{Var}(T) = \Theta(N^2)$$

$$T \approx N \log N \pm N$$

So $P(T > n) \approx 0$ for $n \gg N \log N$

"Pf" (LB):

$$\|P_{\text{Id}}^n - \pi\|_{TV} = \max_{ACS} |P_{\text{Id}}^n(A) - \pi(A)|$$

$$\geq |P_{\text{Id}}^n(A) - \pi(A)| \quad \forall ACS$$

Choose $A_k = \{ \begin{array}{l} k \text{ bottom cards of the deck} \\ \text{are ordered} \end{array} \}$ k fixed
($k=10$)

$$\pi(A_k) = \frac{1}{k!} \sim \text{small}$$

To check: $P_{\text{Id}}^n(A) \sim 1$ if $n \ll N \log N$

Observation: while k ^{different} cards have never been picked, at least k bottom cards of the deck will be ordered (because we started from the Id permutation)

The average to pick $N-k$ different cards is

$$\mathbb{E}(T_k) = \frac{N}{N} + \frac{N}{N-1} + \frac{N}{N-2} + \dots + \frac{N}{k} = N \sum_{e=k}^N \frac{1}{e} \approx N(\log N - \log k)$$
$$\text{Var}(T_k) = \Theta(N^2) = N \log\left(\frac{N}{k}\right)$$

Conclusion:

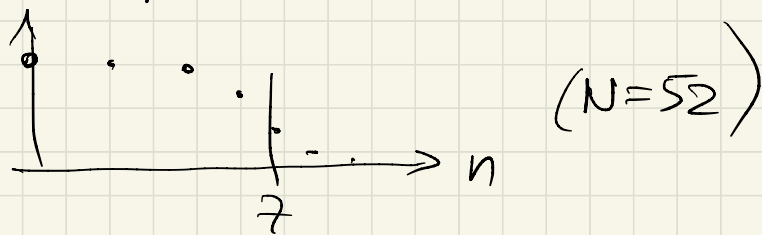
If $n \ll N \log N$, $P_{\text{Id}}^n(A_k) \sim 1$

So $\|P_{\text{Id}}^n - \pi\|_{TV} \sim 1$

Method 2: riffle shuffle

$\Rightarrow \Theta(\log N)$ suffice!

réf: P. Diaconis:



Sampling

$$\pi_i \geq 0, \sum_{i=1} \pi_i = 1$$

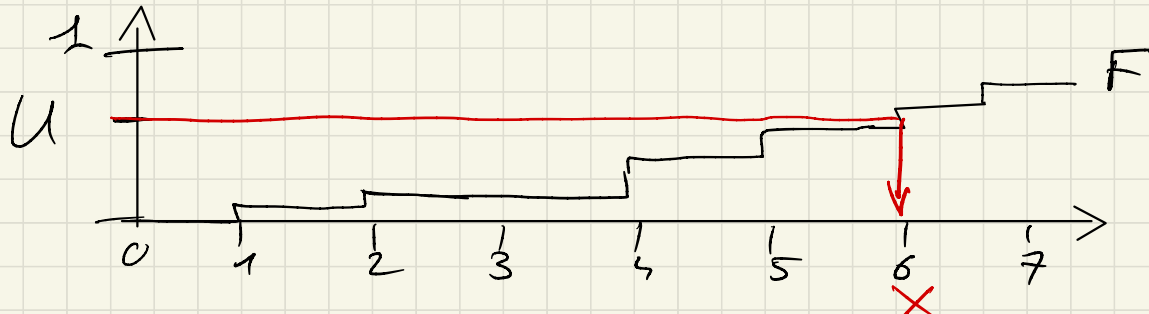
Given a distribution $(\pi_i, i \in S)$ on a state space S , how can we sample from it?

^u Easy solution: let X be a r.v. with values in S such that $P(X=i) = \pi_i$ $i \in S$ and assume $S = \mathbb{N}$

Generate $U \sim \mathcal{U}[0, 1]$ and declare

$$X = \begin{cases} 0 & \text{if } 0 \leq U \leq \pi_0 \\ 1 & \text{if } \pi_0 < U \leq \pi_0 + \pi_1 \\ 2 & \text{if } \pi_0 + \pi_1 < U \leq \pi_0 + \pi_1 + \pi_2 \\ \vdots & \end{cases}$$

$$= F^{-1}(U) \quad \text{where } F = \text{cdf of } X$$



Why to sample? 1. Optimization of a
complex fn

$$f: \{0, 1\}^d \rightarrow \mathbb{R}$$

Am: to maximize f , ie

to find $x_0 \in \{0, 1\}^d$ st $f(x_0) = \max$.

$$\text{Define } \pi(x) = \begin{cases} 0 & \text{if } f(x) \neq \max \\ C = \frac{1}{Z} & \text{if } f(x) = \max \end{cases}$$

where $Z =$ number of maxima of f

First idea: sample from π to get a maximum x

Second idea: instead of sampling from π ,

Sample from π_β defined as follows:

$$\pi_\beta(x) = \frac{e^{\beta f(x)}}{Z_\beta} \quad x \in \{0, 1\}^d$$

where $Z_\beta = \sum_{x \in \{0, 1\}^d} e^{\beta f(x)}$ normalization cst
"partition function"

2. Compute averages (Monte Carlo method)

- $X = \text{r.v.}$ with values in S and distribution π
ie. $\mathbb{P}(X=i) = \pi_i \quad i \in S$
- $f: S \rightarrow \mathbb{R}$
- Aim: Compute $\mathbb{E}(f(X)) = \sum_{i \in S} f(i) \pi_i$

MC method: draw independent samples x_1, \dots, x_M
& compute $\frac{1}{M} \sum_{j=1}^M f(x_j) \sim \mathbb{E}(f(X))$

$$\begin{aligned}\text{Var} \left(\frac{1}{M} \sum_{j=1}^M f(x_j) \right) &= \frac{1}{M^2} \sum_{j=1}^M \underbrace{\text{Var}(f(x_j))}_{= \text{Var}(f(x))} \\ &= \frac{\text{Var}(f(x))}{M}\end{aligned}$$

$$\text{std dev} \sim \Theta \left(\frac{1}{\sqrt{M}} \right)$$