

EPFL

Dependability through Redundancy

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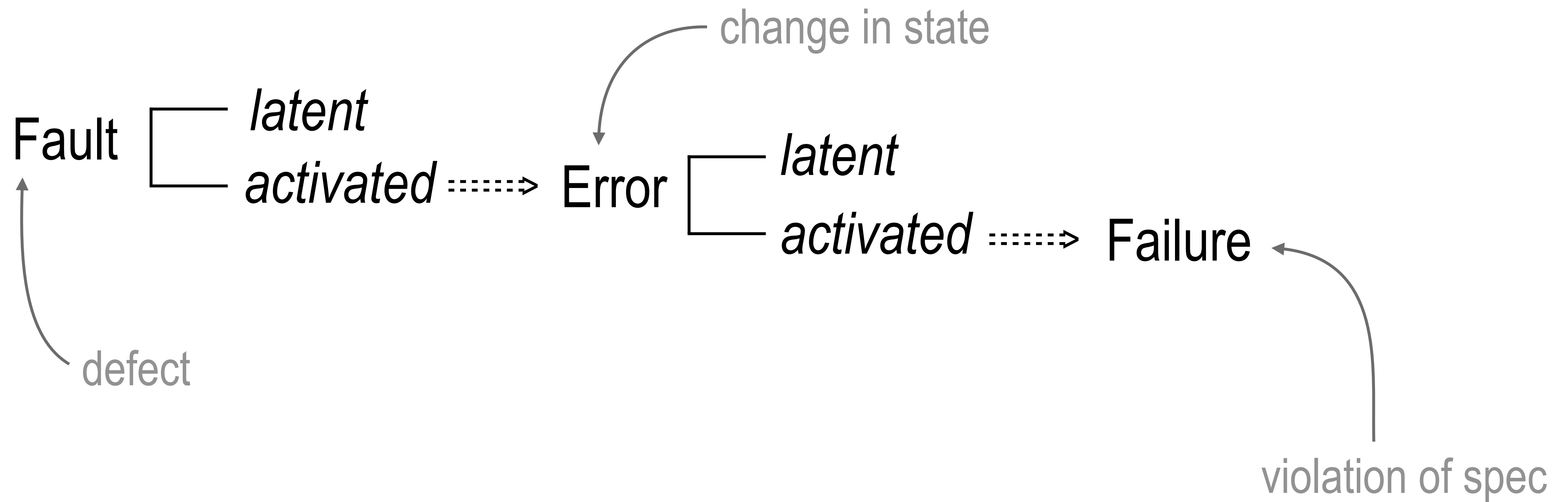
School of Computer & Communication Sciences

How to achieve dependability?

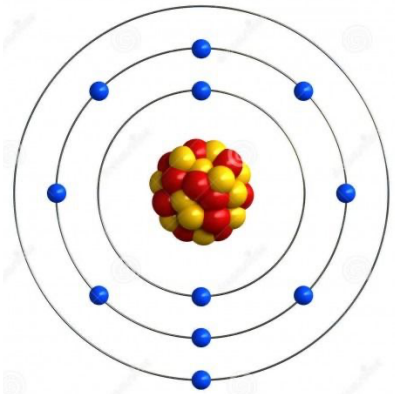
- Use modularity ...
- ... and REDUNDANCY for ...
 - *fault tolerance*
 - *high reliability*
 - *high availability*

Redundancy = duplication with the purpose of increasing dependability

Fault tolerance



Types of software faults / defects



- Bohrbug
 - *clear + easy to reproduce => easy to fix*

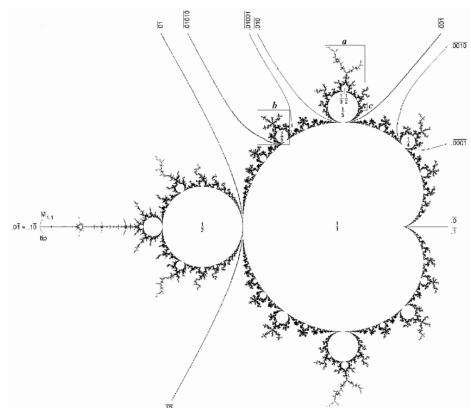
$$\Delta x \Delta p \geq \frac{\hbar}{2}$$

- Heisenbug
 - *disappears when you attach with debugger*



- Schrödingbug
 - *starts causing failure once you realize it should*

- Mandelbug
 - *complex, obscure, chaotic, seemingly non-deterministic*



Using redundancy to tolerate faults

- "tolerate" faults = cope with errors or the resulting failures
- *the actual goal is to tolerate the consequences of faults*

- Using redundancy to cope with errors

Data/information redundancy

- *forward error correction*
- *redundant copies/replicas (=coarse-grained ECC)*
- ...

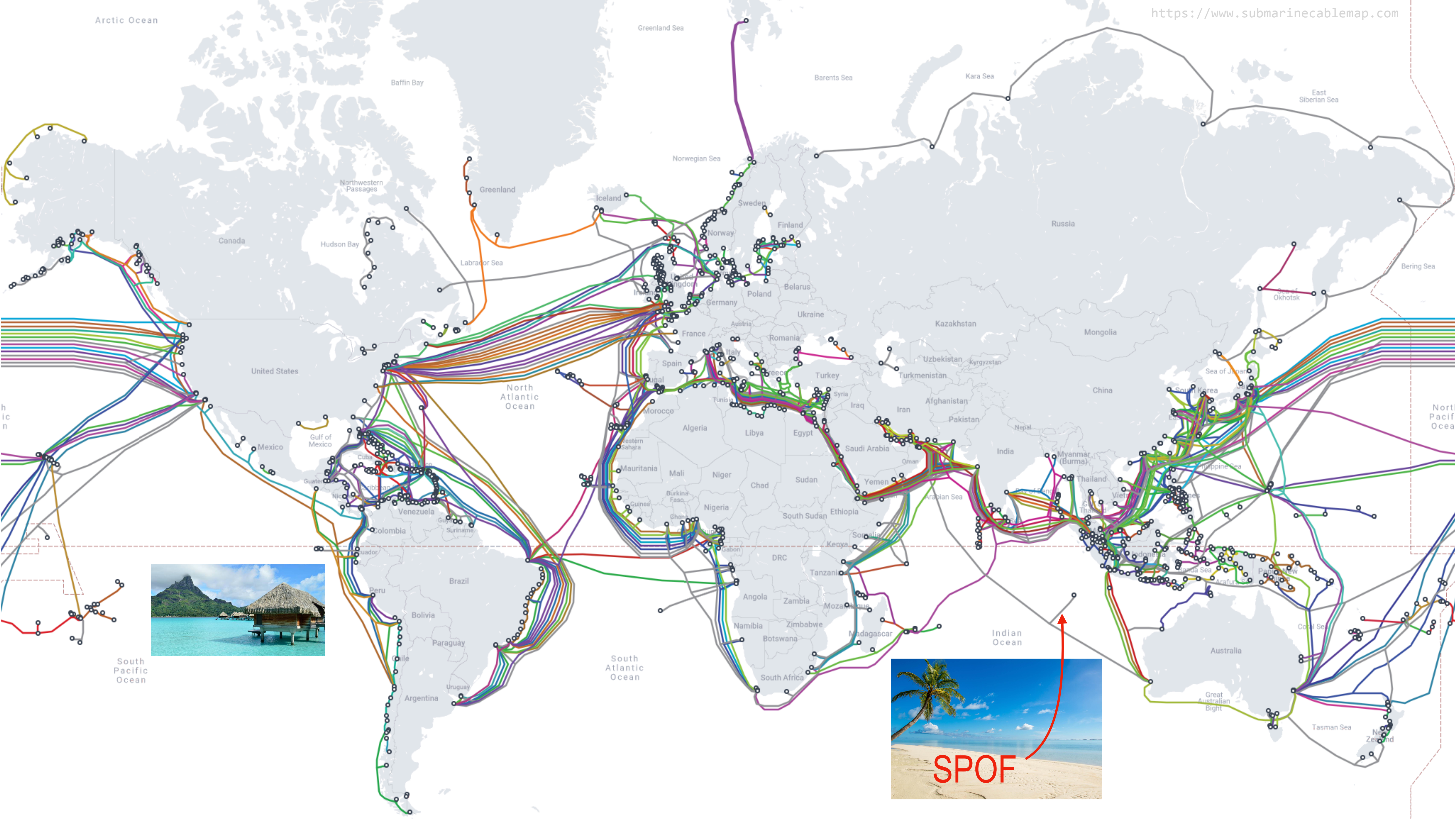
Geographic redundancy

- Using redundancy to cope with failures

Processing redundancy { Space
Time

- *server/service failover*
- *Internet routing*
- ...

Functional redundancy



South Pacific Ocean

South Atlantic Ocean

Indian Ocean

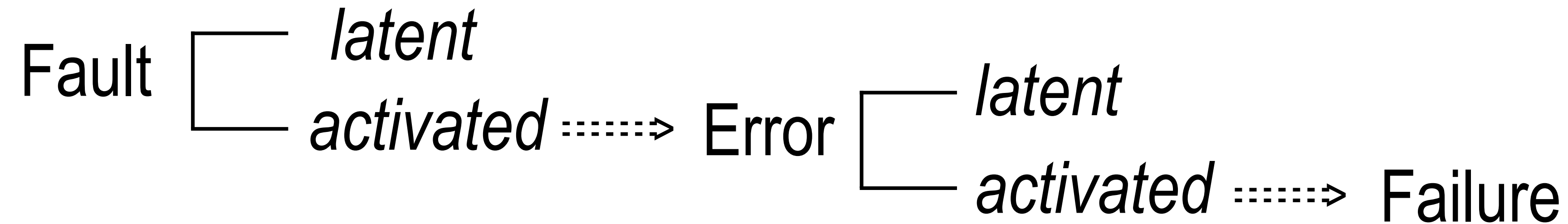
Australia

New Zealand

Fault model

- Specification of what could go wrong and what cannot go wrong
 - *Used to predict consequences of failures*
 - *Should also specify what can / cannot happen during recovery*
 - *Remember the single points of failure (SPOFs)*
- Example: N-version programming
 - *use redundancy to tolerate software faults*

Recap: Fault tolerance



- Different types of software defects
 - *Bohrbug, Heisenbug, ...*
- Redundancy helps tolerate errors and failures
 - *Data redundancy, processing redundancy, ...*
- Fault model = assumptions about what can vs. cannot go wrong

Safety-critical systems

- Safety critical = system whose failure may result in "bad" outcomes
 - *SCADA, aviation, space, automotive, healthcare, ...*
- Fail-safe = failure does not have "bad" consequences
 - *safety-critical \Rightarrow fail-safe*

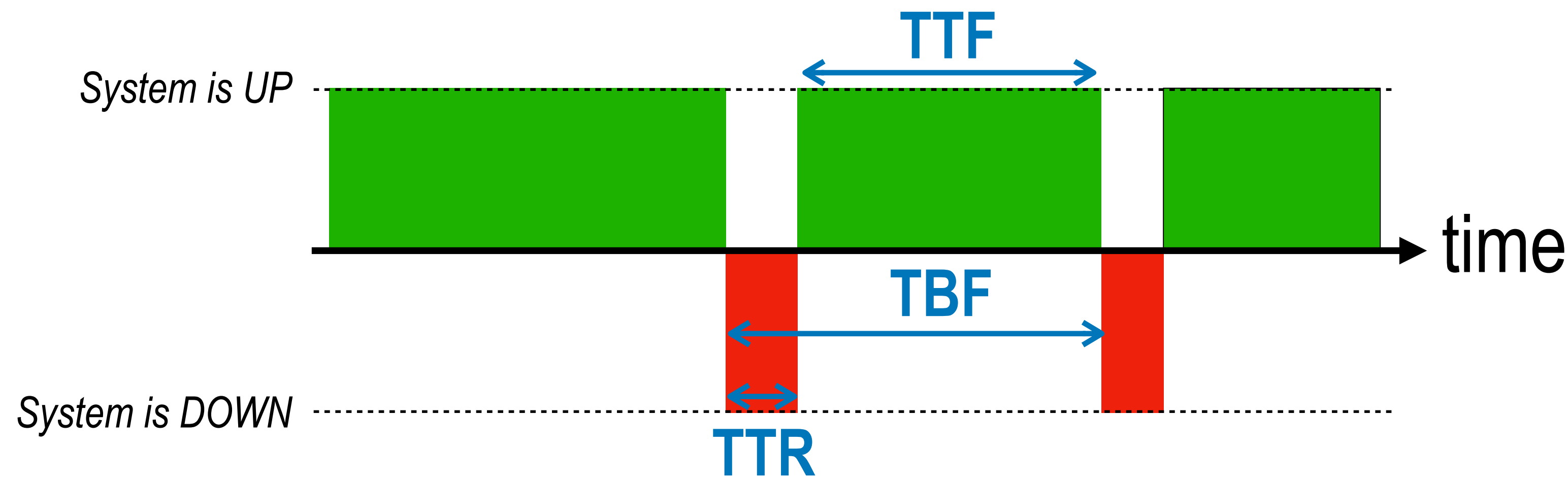
Dependable systems

- Availability = readiness for correct service
- Reliability = continuity of correct service
- Safety = absence of catastrophic consequences
- Confidentiality = absence of unauthorized disclosure of information
- Integrity = absence of improper system state alterations
- Maintainability = ability to undergo repairs and modifications

Reliability

- Reliability = probability of continuous operation
- *continuous operation = (correctly) producing outputs in response to inputs*

$R(t) = P(\text{module operates correctly at time } t \mid \text{it was operating correctly at } t=0)$



$$\text{MTBF} = \text{MTTF} + \text{MTTR}$$

Measuring reliability

- In general MTBF or MTTF ($MTBF = MTTF + MTTR$)
 - *Specifics: Example from SSD spec sheet: P/E cycles, TBW, GB/day, DWPD, MTBF ...*
 - Example: Samsung SSD 850 Pro SATA
 - *Warranty period = 10 years*
 - *MTBF = 2M hours (228 years)*
 - assumes operation of 8 hrs/day
 - 2.5K SSDs => you'd experience 1 failure every ~100 days ($2M / 8 / 2500$)
- Why different???*

Recap: Reliability

- Dependability = Reliability + Availability + Safety + ...
- Safety-critical vs. reliable
- $MTBF = MTTF + MTTR$

Availability

- Availability = probability of producing (correct) outputs in response to inputs

Level of Availability	Percent of Uptime	Downtime per Year	Downtime per Day
1 Nine	90%	36.5 days	2.4 hrs.
2 Nines	99%	3.65 days	14 min.
3 Nines	99.9%	8.76 hrs.	86 sec.
4 Nines	99.99%	52.6 min. $\uparrow \times 10$.86 sec.
5 Nines	99.999%	5.25 min.	.86 sec.
6 Nines	99.9999%	31.5 sec. $\downarrow \div 10$.6 msec

Availability vs. Reliability

- Continuity of service does not matter (unlike reliability)
 - *In theory: uptime is too strict a measure of availability*
 - *In practice: what's the difference?*
- Uptime \Rightarrow availability but Availability $\not\Rightarrow$ uptime
- Examples of ...
 - *Highly available systems with poor reliability (and how is redundancy used)*
...
 - *Highly reliable systems with poor availability (and how is redundancy used)*
...

System availability

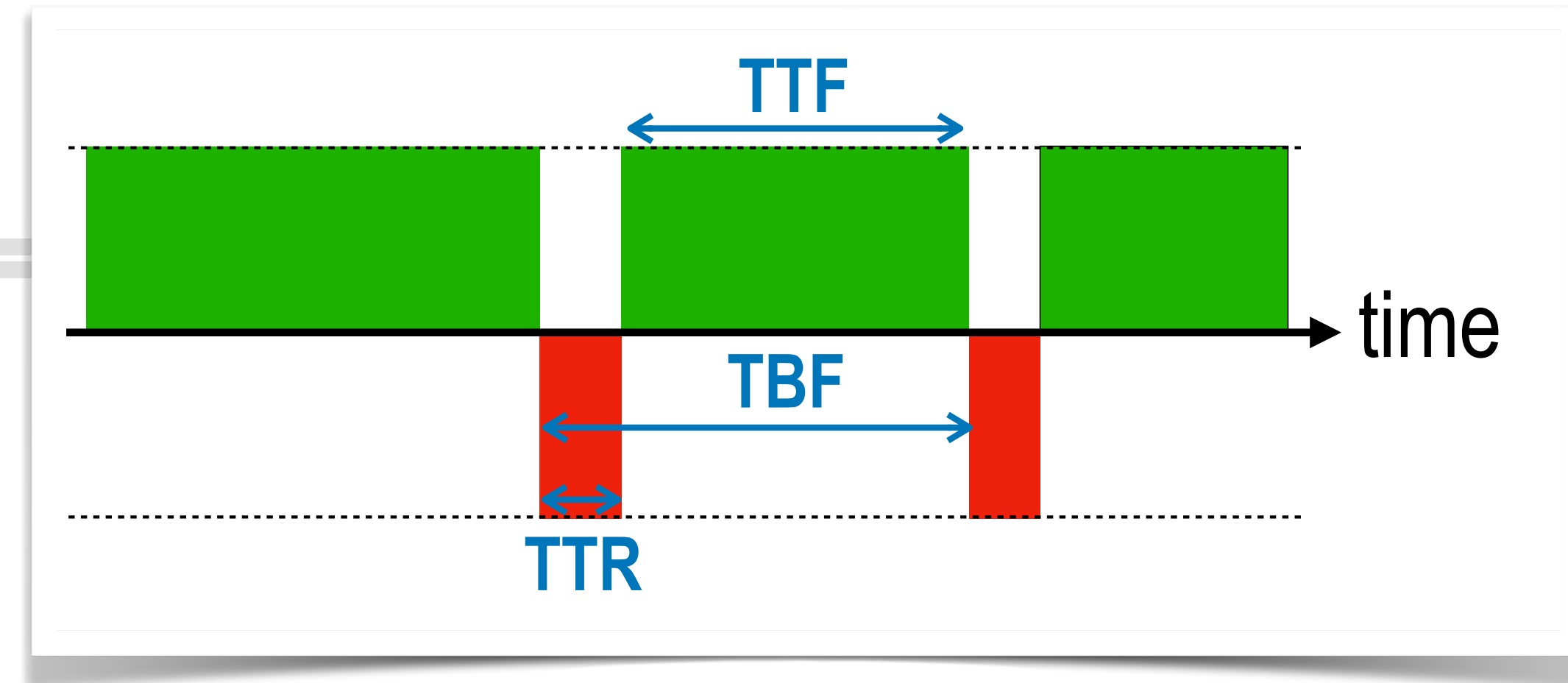
$$\text{Availability} = \frac{\text{MTTF}}{\text{MTBF}}$$

$$\text{Unavailability} = 1 - \text{Availability} = \frac{\text{MTTR}}{\text{MTBF}}$$

$$\text{MTBF} = \text{MTTF} + \text{MTTR} \cong \text{MTTF} \quad (\text{if } \text{MTTF} \gg \text{MTTR})$$

$$\text{Unavailability} \cong \frac{\text{MTTR}}{\text{MTTF}}$$

- Increase availability by
 - *increasing MTTF (higher reliability)*
 - *reducing MTTR (faster recovery)*



Failure modes

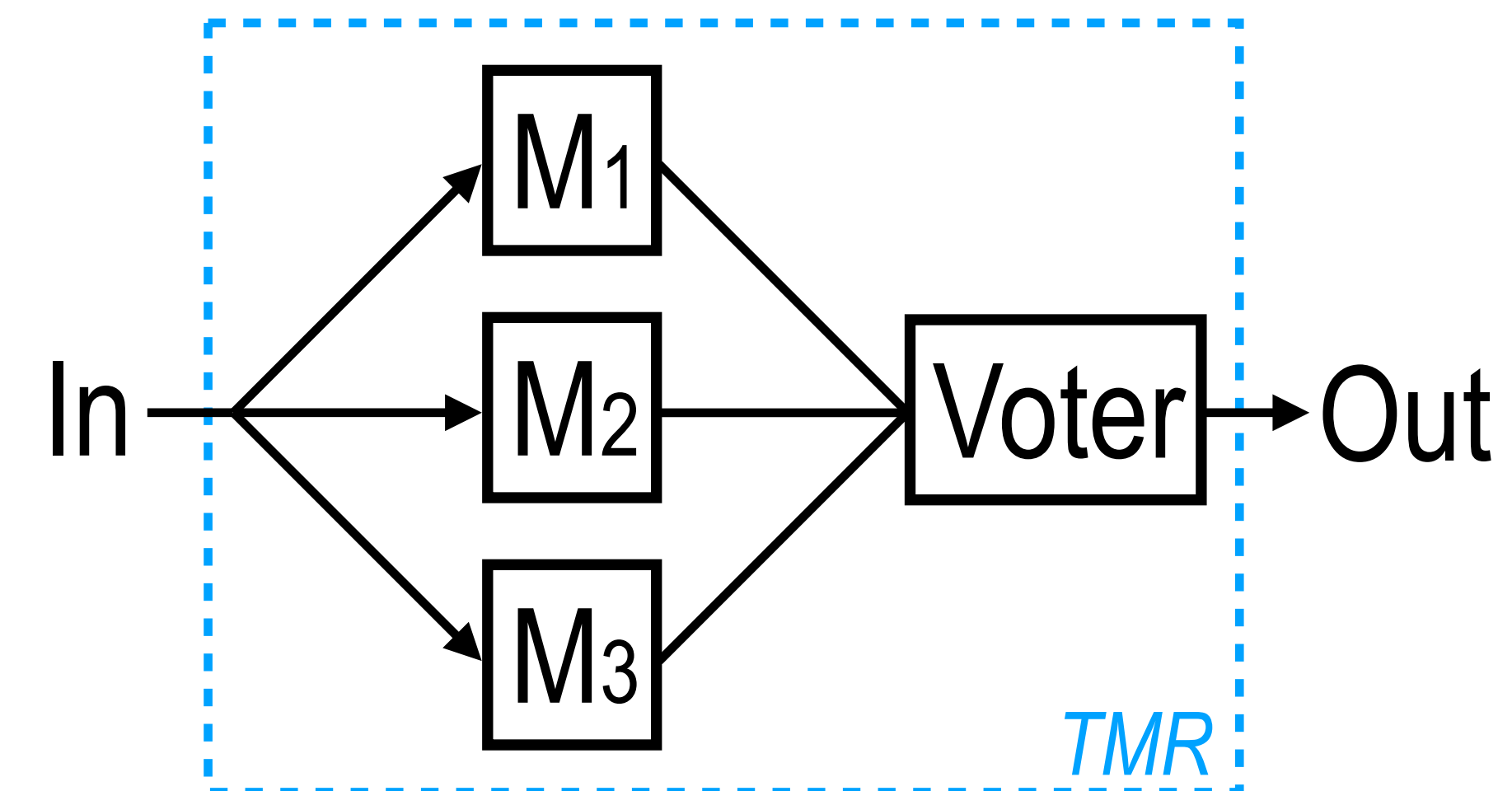
Failure modes

- Definition:

When a system fails, how does that failure appear at the interface of a component?
- Four kinds
 - *fail-stop*
 - *fail-fast*
 - *fail-safe*
 - *fail-soft*

Failure mode 1: Fail-stop

- a.k.a. "crash failure" mode
- *Definition*: halt in response to any internal error that threatens to turn into a failure, before the failure becomes visible
 - \Rightarrow never expose arbitrary behavior
- Any system can be made fail-stop with triple-modular redundancy (TMR)
 - *Strict fault model*: voter is reliable
 - $2f + 1$ independent modules to tolerate f failures
 - *Achilles's heel*: voter



Failure mode 2: Fail-fast

- *Definition:* immediately report at interface any situation that could lead to failure
 - *Can stop immediately after detection or delay (if expect recovery)*
 - *Must stop before failure manifests externally*
- Requires frequent checks of state invariants
- Get auditability of error propagation

Failure mode 3: Fail-safe

- *Definition:* the component remains safe in the face of failure
 - *but possibly degraded functionality or performance*
- "Safety" is context-dependent
- "Controlled" failure

Failure mode 4: Fail-soft

- Definition: internal failures lead to graceful degradation of functionality instead of outright failure

- Example: simple search engine

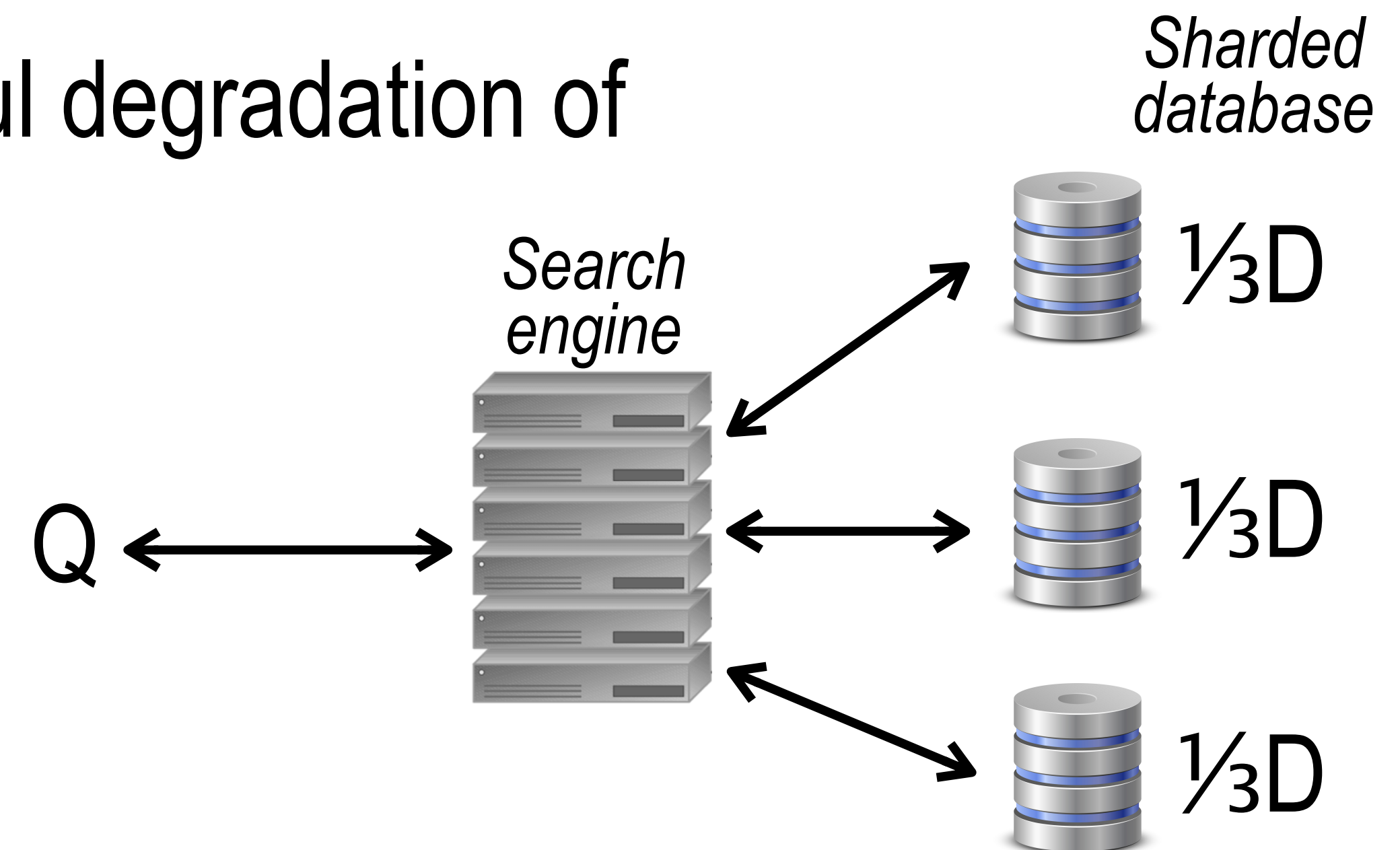
- *system has redundancy at every level*

- Intuition

- *Functionality is typically bottlenecked on data movement (disks, network switches)*

- => Functionality tied to how much data can be moved per unit of time

- **Harvest** (completeness of responses) vs. **yield** (fraction of requests served)



Failure mode 4: Fail-soft: DQ Principle

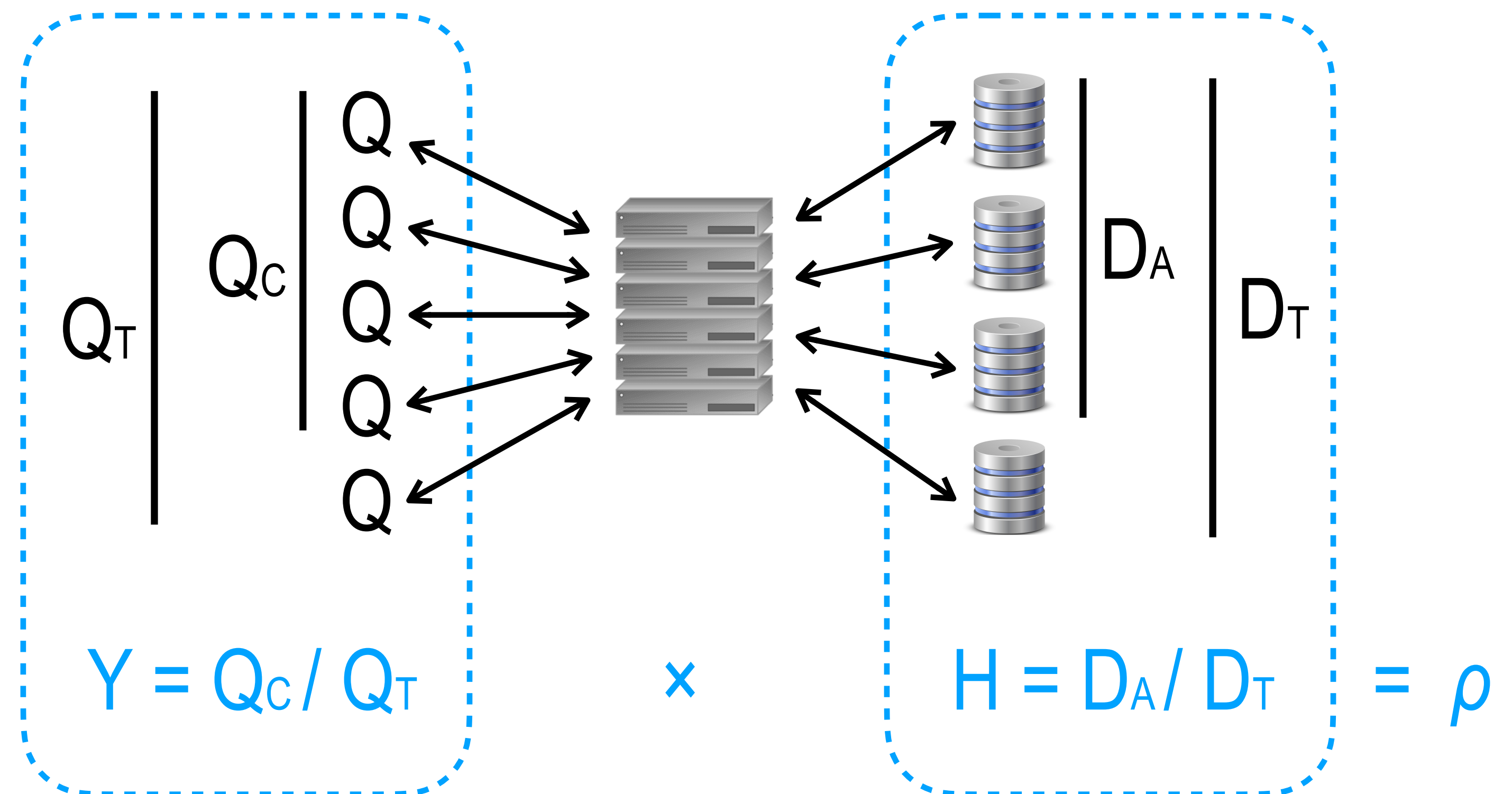
D = data/query
Q = queries/sec

DQ Principle: "D×Q is constant"
(DQ value ρ determined by system configuration)

$$\text{Harvest } H = \frac{D_A}{D_T}$$

$$\text{Yield } Y = \frac{Q_C}{Q_T}$$

DQ Principle: $H \times Y = \rho$ ----->



Recap: Failure modes

- Fail-stop (TMR)
- Fail-fast (Redundant invariant checks)
- Fail-safe
 - *OK to fail, as long as safety is not compromised*
- Fail-soft (Weaker spec)
 - *Redundant resources for top band of acceptable system behavior*
 - *Harvest/yield and the DQ principle in data-intensive parallel systems*

How to improve availability by 10× ?

How to improve availability by 10× ?

$$\text{Unavailability} \cong \frac{\text{MTTR} \downarrow \div 10}{\text{MTTF} \uparrow \times 10}$$

Components of recovery time

- $T_{\text{recover}} = T_{\text{detect}} + T_{\text{diagnose}} + T_{\text{repair}}$
- How to reduce T_{detect} ?
 - Automation
 - Prediction/anticipation
 - Trade-offs between FPs and FNs
- How to reduce T_{diagnose} ?
 - Lots of instrumentation, ML, ...
 - Also a function of what recovery mechanism have available
- How to reduce T_{repair} ?
 - Mostly app-specific
 - Reboot is universal

Detection/Prediction says...

		<i>Detection/Prediction says...</i>	
		Failure	No Failure
<i>Truth is...</i>	No Failure	FP	TN
	Failure	TP	FN

How to improve availability by 10× ?

Reboot-based recovery

Reboot-based recovery

- Design system (components) that recover(s) solely via (micro)rebooting
 - *stop == crash start == recover*
- Design for e.g. microservices
 - *short-running tasks, clusters of many nodes, ...*
- Crash-only components
 - *State segregation*
- Crash-only system of components
 - *Modularization + functional decoupling*
 - *Retryable interactions*
 - *Leased resources*



Reboot-based recovery: State segregation

- Goal: prevent microreboot from inducing corruption or state inconsistency
 - *apply modularization idea to all state: session state vs. persistent state*
- Segment the state by lifetime
- Keep all state that should survive a reboot in dedicated state stores
 - *stores located outside the application ...*
 - *... behind strongly-enforced high-level APIs (e.g., DBs, KV stores)*
- Separate data recovery from app recovery => do each one better

State segregation

Modularization

Functional decoupling

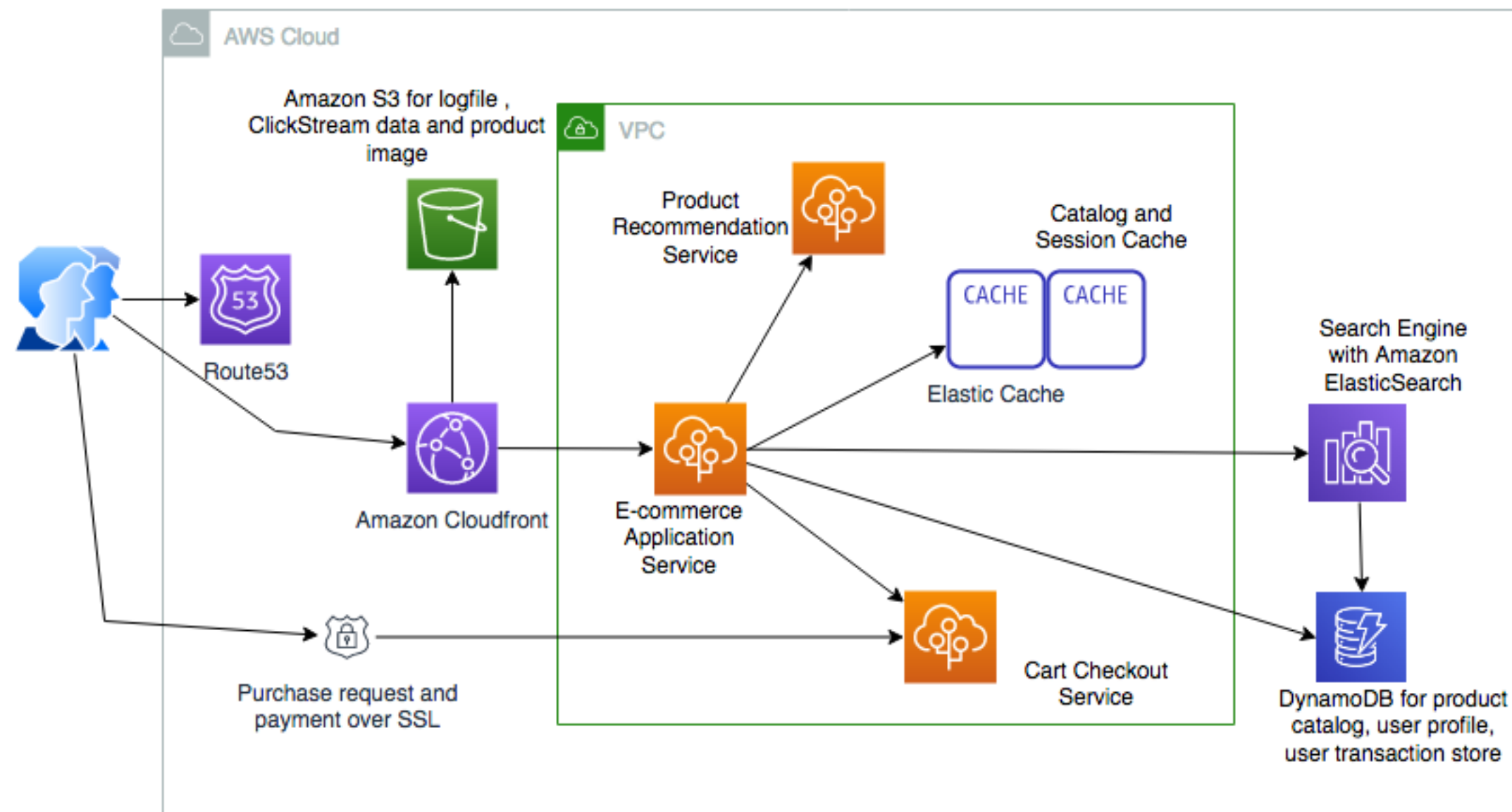
Retryable interactions

Leased resources

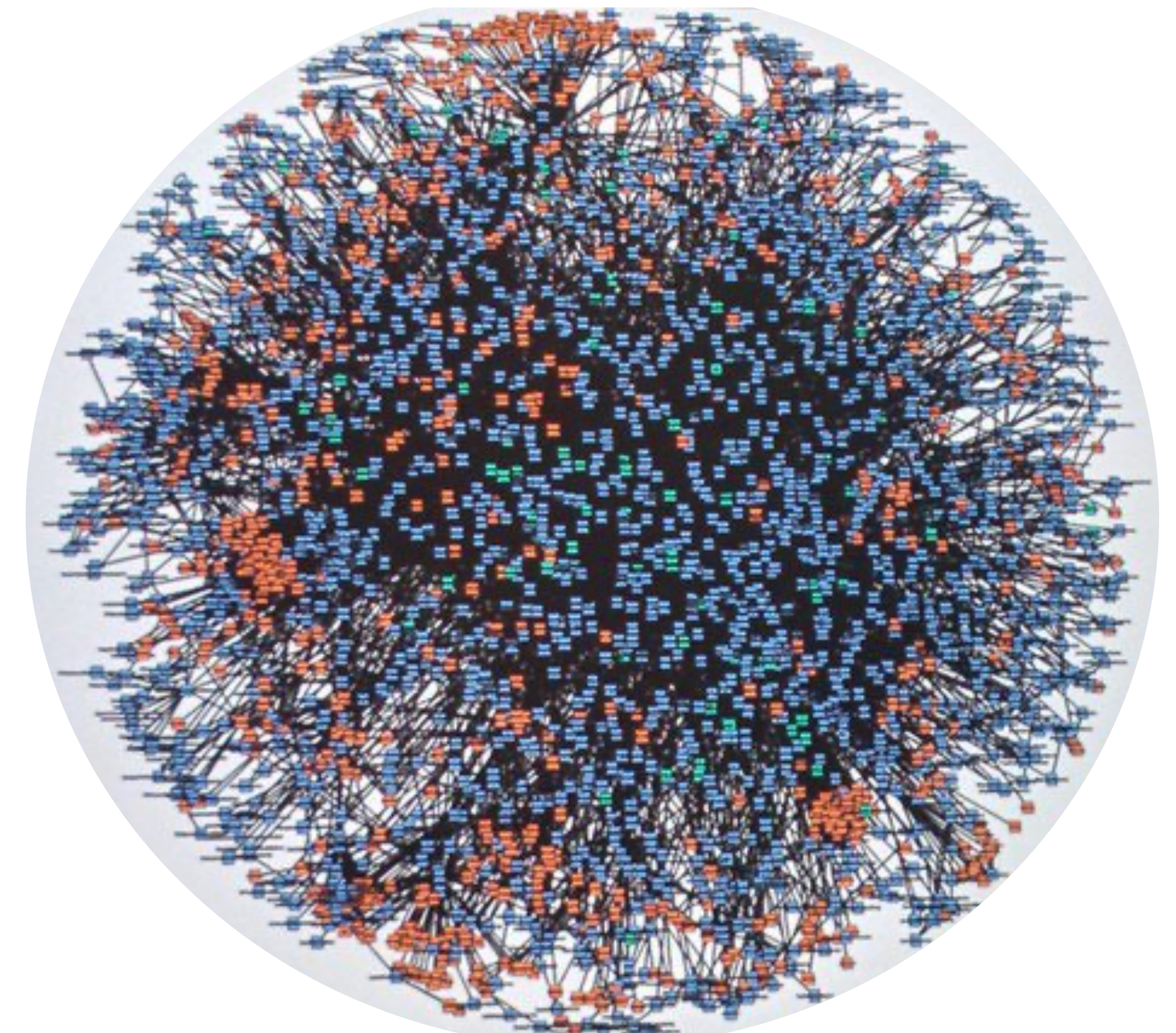
Reboot-based recovery: Strong modularization

- Components with individual loci of control
- *Well defined interfaces*
- *Small in terms of program logic and startup time*
- $T_{\text{reboot}} = T_{\text{restart}} + T_{\text{initialization}}$

State segregation
Modularization
Functional decoupling
Retryable interactions
Leased resources



<https://subscription.packtpub.com/book/web-development/9781838645649/6/ch06/v1/sec28/building-an-soa-based-e-commerce-website-architecture>

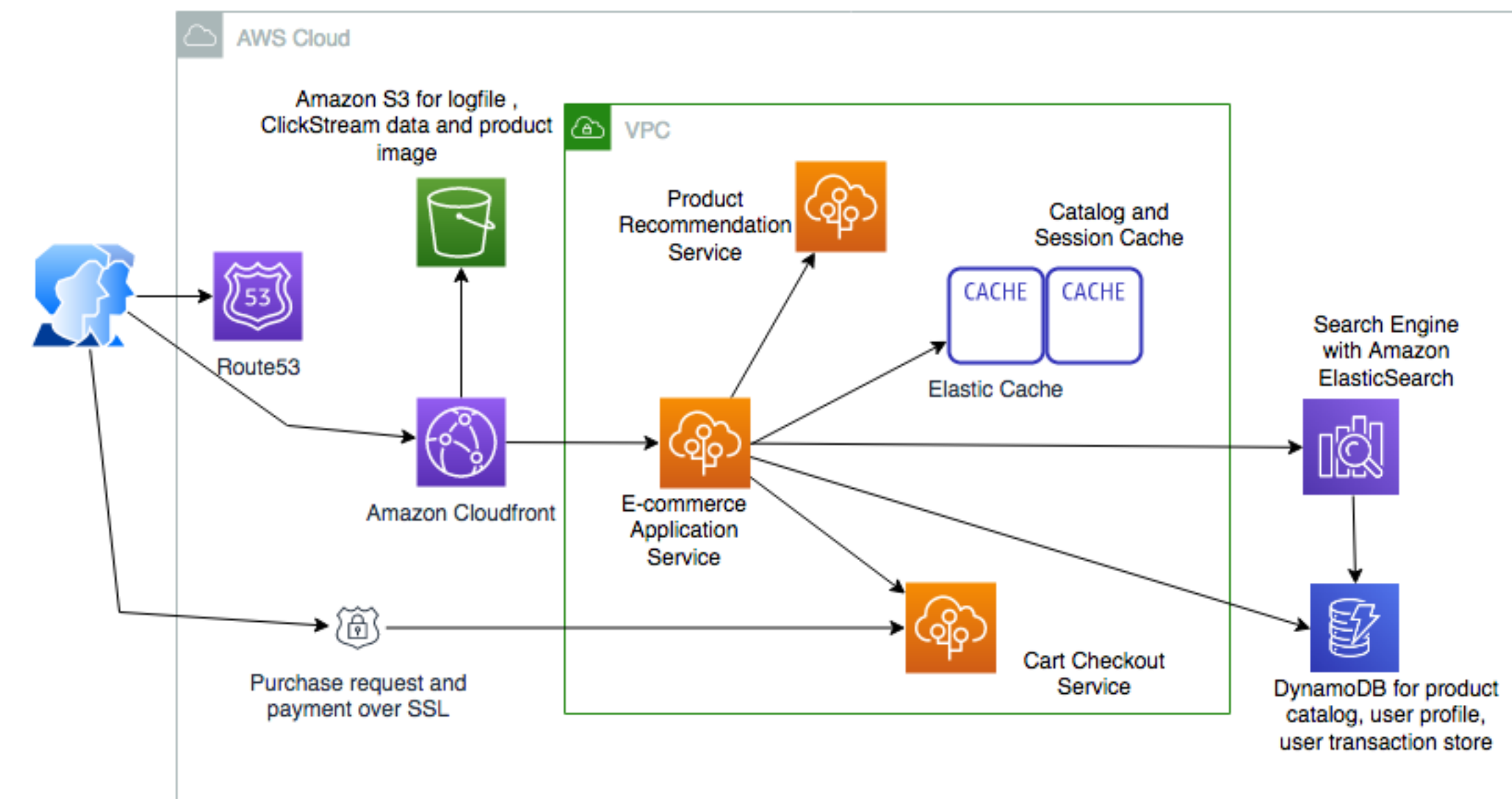


<https://twitter.com/Werner/status/741673514567143424/photo/1>

Reboot-based recovery: Functional decoupling

- Goal
 - *reduced disruption of system during restart*
 - *easy reintegration of component after reinit*
- No direct references (e.g., no pointers) across component boundaries
- *Store cross-component references outside component*
 - Naming indirection through runtime
 - Marshall names into state store

State segregation
Modularization
Functional decoupling
Retryable interactions
Leased resources



<https://subscription.packtpub.com/book/web-development/9781838645649/6/ch06iv1sec28/building-an-soa-based-e-commerce-website-architecture>

Reboot-based recovery: Retryable interactions

- Goal
 - *seamless reintegration of microrebooted component by recovering in-flight requests transparently*
- Interact via timed RPCs or equivalent
 - *if no response, caller can gracefully recover*
 - *timeouts help turn non-Byzantine failures into fail-stop events*
 - *RPC to a microrebooting module throws `RetryAfter(t)` exception*
- Action depends on whether RPC is idempotent or not

State segregation
Modularization
Functional decoupling
Retryable interactions
Leased resources

Exercise: Reboot-based recovery: Leased resources

- Goal: avoid resource leakage without fancy resource tracking
- Lease = timed ownership
 - *File descriptors, memory, ...*
 - *Persistent long-term state*
 - *CPU execution time*
- Requests carry TTL => automatically purged when TTL runs out

State segregation
Modularization
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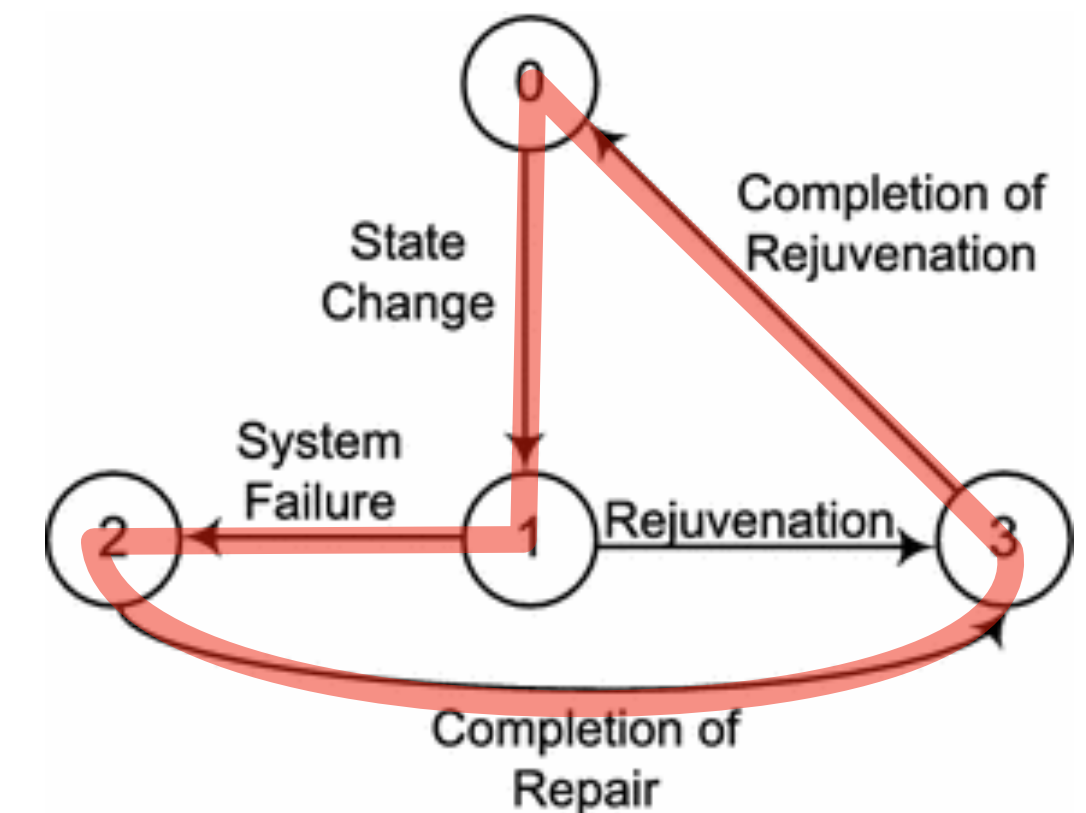
Recap

- $T_{\text{recover}} = T_{\text{detect}} + T_{\text{diagnose}} + T_{\text{repair}}$
- *If recovery is cheap (i.e., T_{repair} is small), can tolerate FPs*
- *Instead of trying to increase MTTF, consider reducing MTTR*
 - Availability goes up, reliability is not affected (in a well designed system)
- **Reboot as a universal "hammer" for curing failures**
 - Systematically employ rebooting to cure failures?
- **Well suited for workloads consisting of fine-grained requests**
 - Currently used in Internet services/microservices, analytics engine, satellite ground station
 - If a fine-grained microreboot doesn't make the problem go away, try coarser-grained

Google "crash-only software" for more info...

Software rejuvenation

- Goal: clean up state to prevent accumulation of errors
 - *Insight: Reboot as a prophylactic*
 - *Does nothing about defects, but reduces probability of turning errors into failures*
- Turns unplanned downtime into planned downtime
 - *Dynamic version of "preventive maintenance"*
 - *Release leaked resources, wipe out data corruption, ...*



Software rejuvenation

- Goal: clean up state to prevent accumulation of errors
 - *Insight: Reboot as a prophylactic*
 - *Does nothing about defects, but reduces probability of turning errors into failures*
- Turns unplanned downtime into planned downtime
 - *Dynamic version of "preventive maintenance"*
 - *Release leaked resources, wipe out data corruption, ...*
- Microrejuvenation
 - *turn unplanned downtime into planned partial downtime (or none at all)*

