

Quantum computation: lecture 9

Shor's algorithm: conclusion

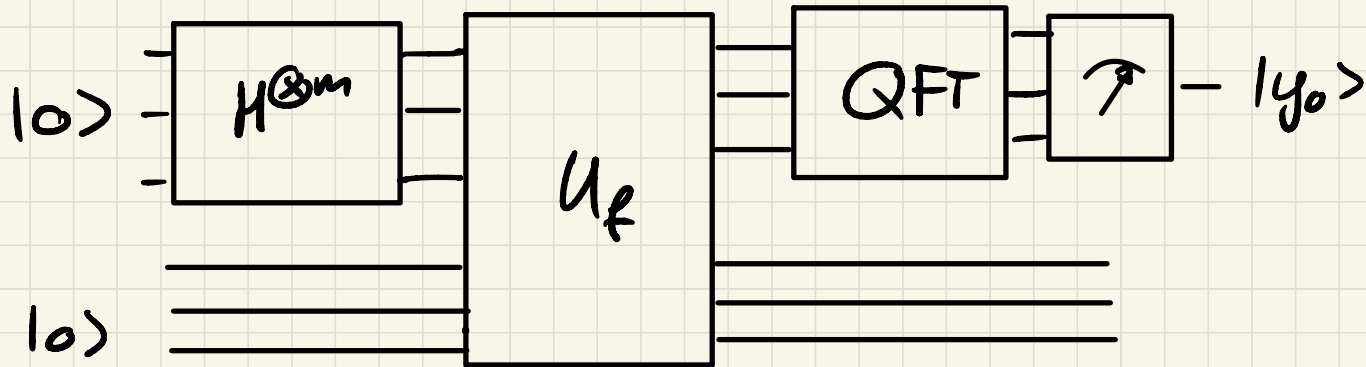
Reminder:

We are looking for the period $r \in \{1..N-1\}$ of a function $f: \mathbb{Z} \rightarrow \mathbb{Z}$ defined as

$$f(x) = a^x \bmod N$$

where $a \in \{1..N-1\}$ is some fixed integer.

For this, we take $M = 2^m$ for some integer $m \geq 1$ such that $M \geq N^2$ (justification later) and use Shor's quantum circuit with $2m$ qubits:



As seen last week, the output of the circuit is a number $y \in \{0..M-1\}$ such that

$$\mathbb{P}(y \in I) \geq \frac{2}{5}$$

where $I = \bigcup_{k=0}^{r-1} I_k$, $I_k = [k \cdot \frac{M}{r} - \frac{1}{2}, k \cdot \frac{M}{r} + \frac{1}{2}]$

i.e. $\exists 0 \leq k \leq r-1$ st. $|y - k \cdot \frac{M}{r}| \leq \frac{1}{2}$

let us divide by M : $|\frac{y}{M} - \frac{k}{r}| \leq \frac{1}{2M}$

Here, the choice of $M \sim N^2 \geq r^2$ matters,
as this implies: $\left| \frac{y}{M} - \frac{k}{r} \right| \leq \frac{1}{2r^2}$

Task: Find in an effective manner all
rational approximations of the form $\frac{k}{r}$
that are at most $\frac{1}{2r^2}$ away from
the measured value $\frac{y}{M}$.

Parenthesis: continued fractions

Pick a real number, for example $\frac{263}{189}$

One-digit approximation: 1

$$\text{So } \frac{263}{189} = 1 + \frac{74}{189} = 1 + \frac{1}{189/74}$$

$$\frac{189}{74} = 2 + \frac{41}{74} = 2 + \frac{1}{74/41}$$

$$\text{So } \frac{263}{189} = 1 + \frac{1}{2 + \frac{1}{74/41}}$$

This leads finally to $\frac{263}{189} = 1 + \frac{1}{2 + \frac{1}{1 + \frac{1}{1 + \frac{1}{4 + \frac{1}{8}}}}}$

$\left(\frac{74}{41} = 1 + \frac{1}{41/33}, \frac{47}{33} = 1 + \frac{1}{33/8}, \frac{33}{8} = 4 + \frac{1}{8} \right)$

Please note $\frac{8}{1} = 7 + \frac{1}{1}$ so this could go on forever, but we choose the shortest development

Note also that if the initial number is irrational then the development is infinite

Notation: $\frac{263}{189} = [1, 2, 1, 1, 4, 8]$

Convergents (toward $\frac{263}{189} \approx 1,391\dots$)

	<u>value</u>	<u>error</u>
[1]	1	$ \frac{263}{189} - 1 < \frac{1}{2}$
[1, 2]	$1 + \frac{1}{2} = \frac{3}{2} = 1,5$	$ \frac{263}{189} - \frac{3}{2} < \frac{1}{9}$
[1, 2, 1]	$1 + \frac{1}{2 + \frac{1}{1}} = \frac{4}{3} = 1, \bar{3}$	$ \frac{263}{189} - \frac{4}{3} < \frac{1}{18}$
[1, 2, 1, 1]	⋮	⋮
[1, 2, 1, 1, 4]	⋮	⋮
[1, 2, 1, 1, 4, 8]	$\frac{263}{189}$	0

Legendre: Let α be a real number

Let p, q be so that $|\alpha - \frac{p}{q}| < \frac{1}{2q^2}$

Then $\frac{p}{q}$ is a convergent of α .

Algorithm:

- Given y , compute the continued fraction of $\frac{y}{M}$
- Look at all convergents: check if any of the denominators is a valid period. If yes, we are done; if not, try again with another measurement.

Note: As $|\frac{y}{n} - \frac{k}{r}| \leq \frac{1}{2n^2}$, we know by Legendre's lemma that $\frac{k}{r}$ must be a convergent of $\frac{y}{n}$, which justifies the previous algorithm!

Complexity: • computing the convergents of $\frac{y}{n}$ is actually Euclid's algorithm for computing $\text{gcd}(y, n)$: at most $O(\log_2 n) = O(m)$ steps
• each division costs $O(m^2)$
 $\Rightarrow O(m^3)$ complexity in total.

Circuit for the QFT

$$\text{QFT} |x\rangle = \frac{1}{2^{m/2}} \sum_{y=0}^{2^m-1} e^{2\pi i xy/2^m} |y\rangle$$

$$\underline{m=1}: \text{QFT} |x\rangle = \frac{1}{\sqrt{2}} (|0\rangle + e^{\frac{2\pi i x}{2}} |1\rangle)$$

$$= \frac{1}{\sqrt{2}} (|0\rangle + (-1)^x |1\rangle) = H |x\rangle$$

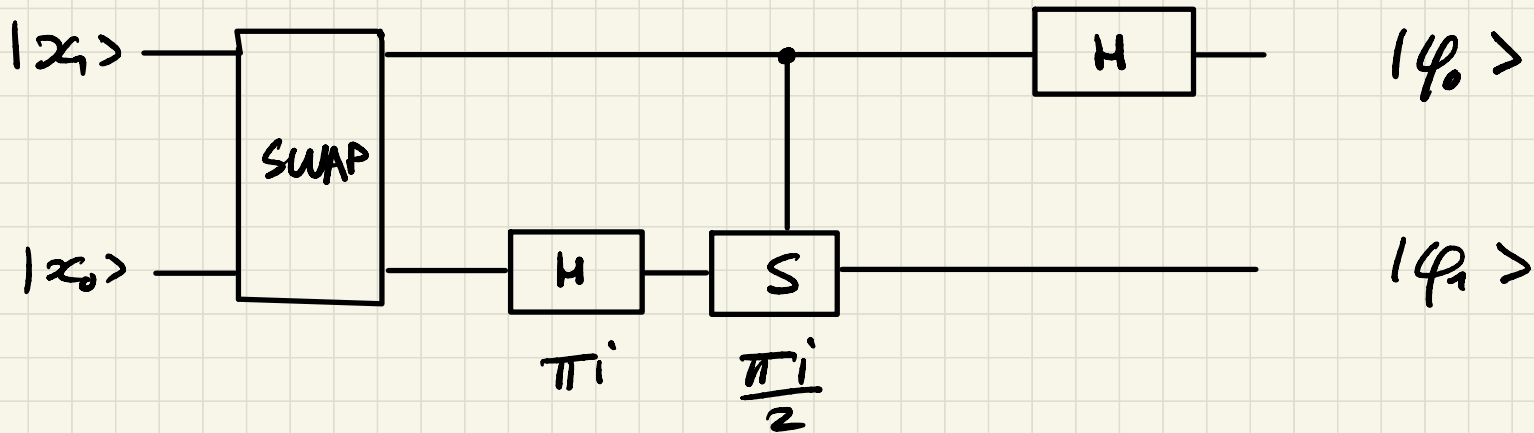
$$\underline{m=2}: \text{QFT} |x\rangle = \frac{1}{2} \left(\overset{00}{\uparrow} |0\rangle + e^{\frac{\pi i x}{2}} \overset{01}{\uparrow} |1\rangle + e^{\pi i x} \overset{10}{\uparrow} |2\rangle + e^{\frac{3\pi i x}{2}} \overset{11}{\uparrow} |3\rangle \right)$$

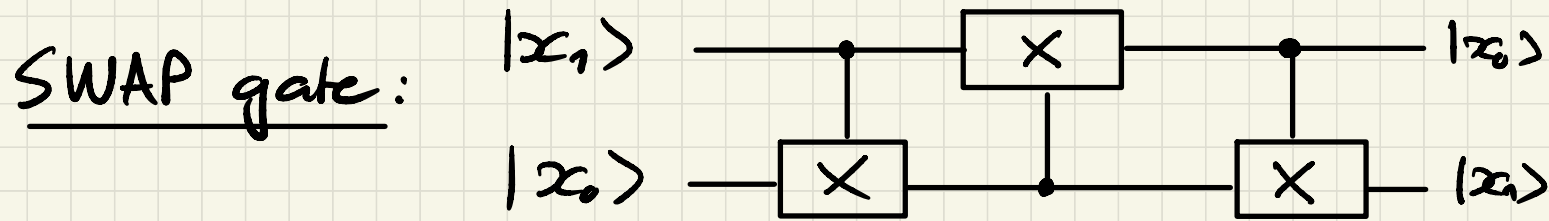
$$= \frac{1}{\sqrt{2}} (|0\rangle + e^{\pi i x} |1\rangle) \otimes \frac{1}{\sqrt{2}} (|0\rangle + e^{\frac{\pi i x}{2}} |1\rangle)$$

Write $x = (x_1, x_0) = 2x_1 + x_0$ (binary expansion)

$$\Rightarrow \text{QFT } |x\rangle = \frac{1}{\sqrt{2}} (|0\rangle + e^{\pi i x_0} |1\rangle) \quad |\varphi_0\rangle$$

$$\otimes \frac{1}{\sqrt{2}} (|0\rangle + e^{\pi i x_1 + \frac{\pi i x_0}{2}} |1\rangle) \quad |\varphi_1\rangle$$





This procedure generalizes to all values of m
 (see next slides)

Circuit complexity:

- $3m$ gates for the swap operations
- $m + (m-1) + (m-2) + \dots + 1 = \frac{m(m+1)}{2} = O(m^2)$ gates for the other part

General m:

Claim: QFT $|x\rangle = \bigotimes_{\ell=1}^m \left(\frac{1}{\sqrt{2}} (|0\rangle + e^{\frac{\pi i x}{2^\ell}} |1\rangle) \right)$

m=1: $\frac{1}{\sqrt{2}} (|0\rangle + e^{\pi i x} |1\rangle)$

m=2: $\frac{1}{\sqrt{2}} (|0\rangle + e^{\pi i x} |1\rangle) \otimes \frac{1}{\sqrt{2}} (|0\rangle + e^{\pi i x/2} |1\rangle)$ ✓

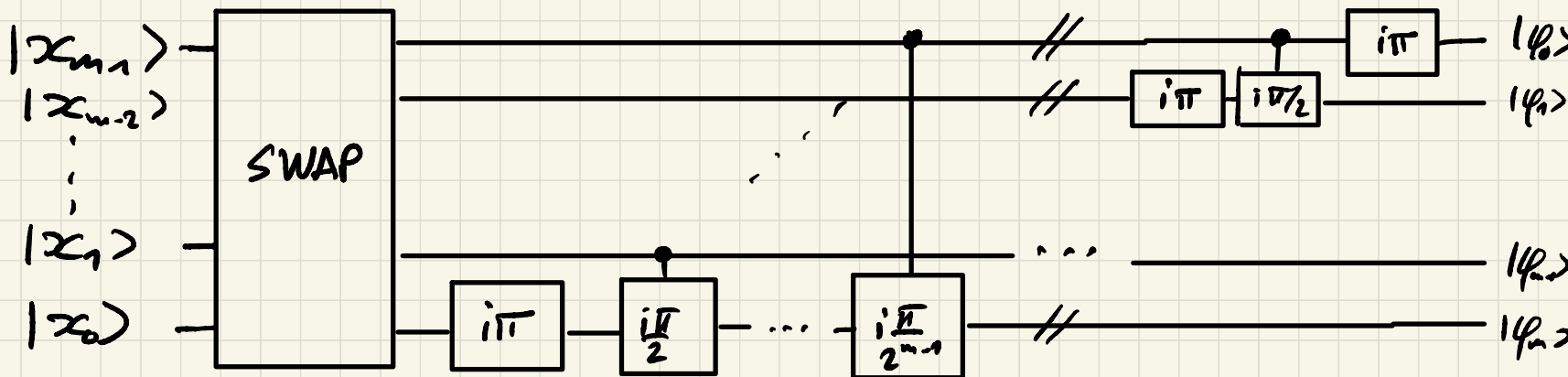
general m:

$$\frac{1}{2^{m/2}} (|0\rangle + e^{\pi i x} |1\rangle) \otimes (|0\rangle + e^{\frac{\pi i x}{2}} |1\rangle) \otimes \dots \otimes (|0\rangle + e^{\frac{\pi i x}{2^{m-1}}} |1\rangle)$$

$$x = (x_{m-1}, \dots, x_0) = x_{m-1} 2^{m-1} + \dots + x_1 2 + x_0 \quad (\text{bin. exp})$$

\Rightarrow QFT $|x\rangle$

$$= \frac{1}{2^{m/2}} \left(|0\rangle + e^{i\pi x_0} |1\rangle \right) \otimes \left(|0\rangle + e^{i\pi x_1 + \frac{\pi i x_0}{2}} |1\rangle \right) \otimes \dots \otimes \left(|0\rangle + e^{i\pi x_{m-1} + \frac{\pi i x_{m-2}}{2} + \dots + \frac{\pi i x_0}{2^{m-1}}} |1\rangle \right)$$



Check of the claim:

$$\text{QFT } |x\rangle = \frac{1}{2^{m/2}} \sum_{y=0}^{2^m-1} e^{\frac{2\pi i x y}{2^m}} |y\rangle$$

For $y = y_0 + 2y_1 + \dots + 2^{m-1}y_{m-1}$, the corresponding phase is: $e^{\pi i x \cdot y_{m-1}} \cdot e^{\frac{\pi i x}{2} \cdot y_{m-2}} \dots e^{\frac{\pi i x}{2^{m-1}} \cdot y_0}$

and one can check for a given sequence of bits y_{m-1}, \dots, y_0 , the phases match in the above expression and that given by the claim. \neq